



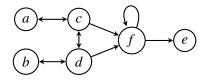
# An Extension-Based Approach to Belief Revision in Abstract Argumentation

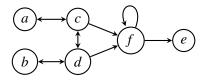
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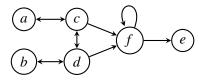
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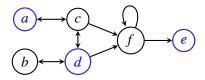


- Evaluation: argumentation semantics
- Extension: set of jointly acceptable arguments



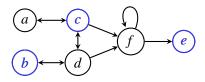
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$$stable(F) =$$



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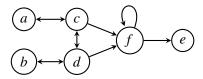
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- Evaluation: argumentation semantics
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$$\textit{stable}(F) = \{\{a, d, e\}, \{b, c, e\}\}$$

Abstract Argumentation Framework (AF) [Dung, 1995]:



- Evaluation: argumentation semantics
- Extension: set of jointly acceptable arguments

$$\textit{stable}(F) = \big\{ \{a, d, e\}, \{b, c, e\} \big\}$$

• Further semantics: preferred, complete, semi-stabe, stage, ...

- (Abstract) argumentation is an inherently dynamic process.
- Revision when new information arises
- Previously: syntax-based revision

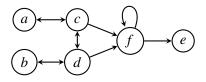
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Model-based revision	Extension-based revision
Knowledge base	Argumentation framework
Model	Extension wrt. $\sigma$
Revision formula	1. Formula / 2. AF
Knowledge base	Argumentation framework

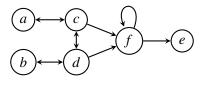
- Coste-Marquis et al., 2014: AGM-style revision of argumentation frameworks, where result is a set of AFs
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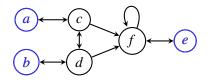
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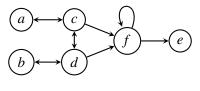
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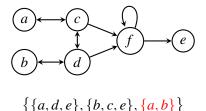
$$\mathit{stable}(F') = \big\{\{a,d,e\},\{b,c,e\},\{a,b,e\}\big\}$$

- Coste-Marquis et al., 2014: AGM-style revision of argumentation frameworks, where result is a set of AFs
- Here: Revision results in a single AF



$$\{\{a,d,e\},\{b,c,e\},\{a,b\}\}$$

- Coste-Marquis et al., 2014: AGM-style revision of argumentation frameworks, where result is a set of AFs
- Here: Revision results in a single AF



There exists no argumentation framework having this extension-set under stable semantics!

### **Main Contributions**

- Representation theorems: Correspondence between revision operators captured by rankings and revision operators given by (extended) set of AGM postulates.
- Revision by propositional formulas

• 
$$\star_{\sigma} : AF_{\mathfrak{A}} \times \mathcal{P}_{\mathfrak{A}} \mapsto AF_{\mathfrak{A}}$$

- Revision by argumentation frameworks
  - $*_{\sigma}: AF_{\mathfrak{A}} \times AF_{\mathfrak{A}} \mapsto AF_{\mathfrak{A}}$

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Revision by argumentation frameworks

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$$*_{\sigma}$$
:  $AF_{\mathfrak{A}} \times AF_{\mathfrak{A}} \mapsto AF_{\mathfrak{A}}$ 

#### Tool-Kit:

- Realizability results for AF semantics [Dunne et al., 2014]
  - Exact characterization of realizable extension-sets  $\Sigma_{\sigma}$
- Horn belief revision [Delgrande and Peppas, 2015]
  - How to modify postulates and rankings in order to stay in the fragment

### **Covered Semantics**

## Definition (Proper I-maximal Semantics)

A semantics  $\sigma$  is called proper I-maximal if for each  $\mathbb{S} \in \Sigma_{\sigma}$ :

- for all  $S_1, S_2 \in \mathbb{S}$ :  $S_1 \subseteq S_2$  implies  $S_1 = S_2$
- ② for all  $\emptyset \neq \mathbb{S}' \subseteq \mathbb{S}$ :  $\mathbb{S}' \in \Sigma_{\sigma}$
- **③** for all ⊆-incomparable extensions  $S_1, S_2$ :  $\{S_1, S_2\} \in \Sigma_{\sigma}$

#### Examples:

- stable semantics
- preferred semantics
- semi-stable semantics
- stage semantics

$$\star_{\sigma} \colon AF_{\mathfrak{A}} \times \mathcal{P}_{\mathfrak{A}} \mapsto AF_{\mathfrak{A}}$$
:

- (P $\star$ 1)  $\sigma(F \star_{\sigma} \varphi) \subseteq [\varphi]$ .
- (P $\star$ 2) If  $\sigma(F) \cap [\varphi] \neq \emptyset$  then  $\sigma(F \star_{\sigma} \varphi) = \sigma(F) \cap [\varphi]$ .
- (P $\star$ 3) If  $[\varphi] \neq \emptyset$  then  $\sigma(F \star_{\sigma} \varphi) \neq \emptyset$ .
- (P\*4) If  $\varphi \equiv \psi$  then  $\sigma(F \star_{\sigma} \varphi) = \sigma(F \star_{\sigma} \psi)$ .
- (P\*5)  $\sigma(F \star_{\sigma} \varphi) \cap [\psi] \subseteq \sigma(F \star_{\sigma} (\varphi \wedge \psi)).$
- $(\text{P} \star 6) \ \text{ If } \sigma(F \star_{\sigma} \varphi) \cap [\psi] \neq \emptyset \text{ then } \sigma(F \star_{\sigma} (\varphi \wedge \psi)) \subseteq \sigma(F \star_{\sigma} \varphi) \cap [\psi].$

[Alchourrón et al., 1985, Katsuno and Mendelzon, 1991, Coste-Marquis et al., 2014]

### Definition ( $\sigma$ -compliance)

A pre-order  $\preceq$  is  $\sigma$ -compliant if for every formula  $\varphi$  it holds that  $\min([\varphi], \preceq)$  is realizable under  $\sigma$ .

## Example ( $\sigma \in \{stable, preferred, stage, semi-stable\}$ )

- $\varphi = \neg(a \land b \land c)$
- $\bullet \{a,b,c\} \prec \{a,b\} \approx \{a,c\} \approx \{b,c\} \prec \{a\} \approx \{b\} \approx \{c\} \prec \emptyset$ 
  - $\min([\varphi], \preceq) = \{\{a, b\}, \{a, c\}, \{b, c\}\} \notin \Sigma_{\sigma}$
  - $\leq$  is not  $\sigma$ -compliant
- $\bullet \ \{a,b,c\} \prec' \{a\} \approx' \{b\} \approx' \{c\} \prec' \{a,b\} \prec' \{a,c\} \prec' \{b,c\} \prec' \emptyset$ 
  - $\leq'$  is  $\sigma$ -compliant
  - For instance,  $\min([\varphi], \preceq') = \{\{a\}, \{b\}, \{c\}\} \in \Sigma_{\sigma}$

#### Definition

Given semantics  $\sigma$  and AF F, a pre-order  $\leq_F$  is a faithful ranking if it is total and for any sets  $E_1, E_2$  and AFs  $F, F_1, F_2$ :

- (i) if  $E_1, E_2 \in \sigma(F)$ , then  $E_1 \approx_F E_2$ ,
- (ii) if  $E_1 \in \sigma(F)$  and  $E_2 \notin \sigma(F)$ , then  $E_1 \prec_F E_2$ ,
- (iii) if  $\sigma(F_1) = \sigma(F_2)$ , then  $\leq_{F_1} = \leq_{F_2}$ .

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#### **Theorem**

An operator  $\star_{\sigma}$  satisfies postulates  $P\star 1 - P\star 6$  for proper I-maximal semantics  $\sigma$ 

iff

there exists an assignment mapping each AF F to a faithful and  $\sigma$ -compliant ranking  $\leq_F$  such that  $\sigma(F \star_{\sigma} \varphi) = \min([\varphi], \leq_F)$ .

$$*_{\sigma}: AF_{\mathfrak{A}} \times AF_{\mathfrak{A}} \mapsto AF_{\mathfrak{A}}:$$

- (A\*1)  $\sigma(F *_{\sigma} G) \subseteq \sigma(G)$ .
- (A\*2) If  $\sigma(F) \cap \sigma(G) \neq \emptyset$ , then  $\sigma(F *_{\sigma} G) = \sigma(F) \cap \sigma(G)$ .
- (A\*3) If  $\sigma(G) \neq \emptyset$ , then  $\sigma(F *_{\sigma} G) \neq \emptyset$ .
- (A\*4) If  $\sigma(G) = \sigma(H)$ , then  $\sigma(F *_{\sigma} G) = \sigma(F *_{\sigma} H)$ .
- (A\*5)  $\sigma(F *_{\sigma} G) \cap \sigma(H) \subseteq \sigma(F *_{\sigma} f_{\sigma}(\sigma(G) \cap \sigma(H))).$
- (A\*6) If  $\sigma(F *_{\sigma} G) \cap \sigma(H) \neq \emptyset$ , then  $\sigma(F *_{\sigma} f_{\sigma}(\sigma(G) \cap \sigma(H))) \subseteq \sigma(F *_{\sigma} G) \cap \sigma(H)$ .
- (*Acyc*) If for  $0 \le i \le n$  we have  $\sigma(F *_{\sigma} G_{i+1}) \cap \sigma(G_i) \ne \emptyset$  and  $\sigma(F *_{\sigma} G_0) \cap \sigma(G_n) \ne \emptyset$  then  $\sigma(F *_{\sigma} G_n) \cap \sigma(G_0) \ne \emptyset$ .

#### Definition

Given semantics  $\sigma$  and AF F, a pre-order  $\leq_F$  is an I-faithful ranking if it is I-total and for any  $\subseteq$ -incomparable sets  $E_1, E_2$  and AFs  $F, F_1, F_2$ :

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An operator  $*_\sigma$  satisfies postulates A\*1 – A\*6 + (Acyc) for proper I-maximal semantics  $\sigma$ 

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⇒ standard model-based revision operators (e.g. [Dalal, 1988]) work.

### Conclusion

#### Summary:

- Extension-based revision resulting in a single AF
- Combining recent results in argumentation and belief revision
- Different representation theorems:
  - Revision by propositional formulas
  - Revision by argumentation frameworks

#### Future work:

- Concrete operators
- Other semantics
- Minimal-change criteria for the realizing AFs
- Iterated revision of AFs

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