6.0/4.0 VU Formale Methoden der Informatik (185.291) December 9, 2016

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1.) Consider the following problem:

PROB

INSTANCE: A program Π such that Π takes a string as input, and outputs a string. It is guaranteed that Π terminates on any input string.

QUESTION: Do there exist strings I_1, I_2 such that $\Pi(I_1) = I_2$, i.e., such that the output of Π on the input I_1 is equal to I_2 ?

Prove that the problem \mathbf{PROB} is semi-decidable. For this, describe a procedure that shows the semi-decidability of the problem (i.e. a semi-decision procedure for \mathbf{PROB}) and argue that it is correct.

Note: we consider only strings that are built from symbols 0 and 1. (15 points)

2.) (a) Given the following first-order logic formula ψ :

$$\psi: \left[p(f(x,y),u) \land p(x,z) \right] \to p(f(z,y),u)$$

where f/2 is a binary function symbol and p/2 is a binary predicate symbol. Let T be a theory which forces p/2 to be reflexive, symmetric, and transitive. Additionally, T includes the following axiom related to p and f:

$$\forall x_1, x_2, y_1, y_2 : [p(x_1, x_2) \land p(y_1, y_2)] \rightarrow p(f(x_1, y_1), f(x_2, y_2))$$

Give a detailed proof that ψ is T-valid.

- (b) Consider the clauses C_1, \ldots, C_5 in **dimacs** format (in this order, shown in the box; recall that 0 indicates the end of a clause) which are given as input to a SAT solver. Apply CDCL to solve the CNF using the convention that if a variable is assigned as a decision, then it is assigned 'false'. Further, select variable 3 as the first decision variable that is assigned.
 - Each time when a conflict occurs and after backtracking, draw the implication graph and indicate all UIPs and mark the first UIP. For each UIP, indicate the cut (i.e., a set of edges) and its asserting conflict clause. Learn the asserting conflict clause that corresponds to the first UIP.
 - Is the given CNF satisfiable, unsatisfiable, or valid? Can the empty clause be derived from the given CNF during CDCL? Justify your answers to the above questions.

-1 -2 -5 0
-1 -2 5 0
2 -4 0
13-40
4 0
4 0

(6 points)

(9 points)

- **3.)** Let π be the program $x \coloneqq x y; y \coloneqq x + y; x \coloneqq y x$.
 - (a) Specify a correctness assertion stating that this program swaps that values of the variables x and y. (1 point)

(b)) Prove the correctness assertion using weakest preconditions. (5 points)
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(c) Prove the correctness assertion using strongest postconditions. (9 points)

4.) (a) Show that simulation is a transitive relation, i.e. given any 3 Kripke structures

$$K_1 = \{S_1, I_1, R_1, L_1\}, K_2 = \{S_2, I_2, R_2, L_2\}$$
 and $K_3 = \{S_3, I_3, R_3, L_3\}$

over atomic predicates AP, such that $K_1 \leq K_2$ and $K_2 \leq K_3$, show that $K_1 \leq K_3$. (5 points)

(b) Consider the following Kripke structure M:



For each of the following formulae $\varphi,$

- i. check the respective box if the formula is in CTL, LTL, and/or CTL*, and
- ii. list the states s_i on which the formula φ holds; i.e. for which states s_i do we have $M, s_i \models \varphi$?

φ	CTL	LTL	CTL^*	States s_i
$\mathbf{G}(b)$				
$\mathbf{F}(a)$				
$\mathbf{X}(a)$				
$\mathbf{A}[a \ \mathbf{U} \ c]$				
$\mathbf{EF}(a)$				

(5 points)

(c) The subset sum problem is defined as follows: Given a set of N integers $S = \{i_1, i_2, \dots, i_N\}$, does S have a nonempty subset whose sum is zero?

Write a C program that implements a *guess and check* routine for the subset sum problem and instrument the program with an appropriate CBMC assertion. You may assume the following template:

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int nondet_bool(); // non-deterministically returns 0 or 1 \,
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(5 points)